## Chapter X02

## Machine Constants

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## 1 Scope of the Chapter

This chapter is concerned with parameters which characterise certain aspects of the computing environment in which the NAG Parallel Library is implemented. They relate primarily to floating-point arithmetic, but also to integer arithmetic and the elementary functions. The values of the parameters vary from one implementation of the Library to another, but within the context of a single implementation they are constants.

The parameters are intended for use primarily by other routines in the Library, but users of the Library may sometimes need to refer to them directly.
Each parameter value is returned by a separate Fortran function. Because of the simple nature of the functions, individual routine documents are not provided; the necessary details are given in Section 3 of this Introduction.

## 2 Background to the Problems

### 2.1 Floating-point Arithmetic

### 2.1.1 A model of floating-point arithmetic

In order to characterise the important properties of floating-point arithmetic by means of a small number of parameters, NAG uses a simplified model of floating-point arithmetic. The parameters of the model can be chosen to provide a sufficiently close description of the behaviour of actual implementations of floating-point arithmetic, but not, in general, an exact description; actual implementations may vary in the details of how numbers are represented or arithmetic operations are performed. In particular, denormalized numbers and gradual underflow is not treated by this model.

The model is based on that developed by Brown [1], but differs in some respects. The essential features are summarised here.

The model is characterised by four integer parameters and one logical parameter. The four integer parameters are:
$b$ : the base
$p$ : the precision (i.e., the number of significant base- $b$ digits)
$e_{\text {min }}$ : the minimum exponent
$e_{\text {max }}$ : the maximum exponent
These parameters define a set of numerical values of the form:

$$
f \times b^{e}
$$

where the exponent $e$ must lie in the range $\left[e_{\min }, e_{\max }\right.$ ], and the fraction $f$ (also called the mantissa or significand) lies in the range $[1 / b, 1)$, and may be written:

$$
f=0 . f_{1} f_{2} \ldots f_{p}
$$

Thus $f$ is a $p$-digit fraction to the base $b$; the $f_{i}$ are the base- $b$ digits of the fraction: they are integers in the range 0 to $b-1$, and the leading digit $f_{1}$ must not be zero.

The set of values so defined (together with zero) are called model numbers. For example, if $b=10, p$ $=5, e_{\min }=-99$ and $e_{\max }=+99$, then a typical model number is $0.12345 \times 10^{67}$.

The model numbers must obey certain rules for the computed results of the following basic arithmetic operations: addition, subtraction, multiplication, negation, absolute value, and comparisons. The rules depend on the value of the logical parameter ROUNDS.
If ROUNDS is true, then the computed result must be the nearest model number to the exact result (assuming that overflow or underflow does not occur); if the exact result is midway between two model numbers, then it may be rounded either way.

If ROUNDS is false, then: if the exact result is a model number, the computed result must be equal to the exact result; otherwise, the computed result may be either of the adjacent model numbers on either side of the exact result.

For division and square root, this latter rule is further relaxed (regardless of the value of ROUNDS): the computed result may also be one of the next adjacent model numbers on either side of the permitted values just stated.

On some machines, the full set of representable floating-point numbers conforms to the rules of the model with appropriate values of $b, p, e_{\min }, e_{\max }$ and ROUNDS. For example, for machines supporting IEEE binary double precision arithmetic:

$$
\begin{array}{llr}
b & = & 2 \\
p & = & 53 \\
e_{\min } & = & -1021 \\
e_{\max } & = & 1024
\end{array} \text { and ROUNDS is true. }
$$

For other machines, values of the model parameters must be chosen which define a large subset of the representable numbers; typically it may be necessary to decrease $p$ by 1 (in which case ROUNDS is always set to false), or to increase $e_{\text {min }}$ or decrease $e_{\max }$ by a little bit. There are additional rules to ensure that arithmetic operations on those representable numbers which are not model numbers are consistent with arithmetic on model numbers.
(Note: the model used here differs from that described in Brown [1] in the following respects: square-root is treated, like division, as a weakly supported operator; and the logical parameter ROUNDS has been introduced to take account of machines with good rounding.)

### 2.1.2 Derived parameters of floating-point arithmetic

Most numerical algorithms require access, not to the basic parameters of the model, but to certain derived values, of which the most important are:
the machine precision $\epsilon:=\frac{1}{2} b^{1-p}$ if ROUNDS is true; $=b^{1-p}$ otherwise (but see Note below).
the smallest positive model number: $=b^{e_{\text {min }}-1}$
the largest positive model number: $=\left(1-b^{-p}\right) b^{e_{\max }}$
Note: this value is increased very slightly in some implementations to ensure that the computed result of $1+\epsilon$ or $1-\epsilon$ differs from 1. For example in IEEE binary double precision arithmetic [2] the value is usually set to $2^{-53}+2^{-105}$ or $2^{53}+2^{-63}$.

Two additional derived values are used in the NAG Parallel Library. Their definitions depend not only on the properties of the basic arithmetic operations just considered, but also on properties of some of the elementary functions. We define the safe range parameter to be the smallest positive model number $z$ such that for any $x$ in the range $[z, 1 / z]$ the following can be computed without undue loss of accuracy, overflow, underflow or other error:

```
\(-x\)
\(1 / x\)
\(-1 / x\)
\(\operatorname{SQRT}(x)\)
LOG \((x)\)
\(\operatorname{EXP}(\operatorname{LOG}(x))\)
\(y * *(\operatorname{LOG}(x) / \operatorname{LOG}(y))\) for any real \(y\)
```

In a similar fashion we define the safe range parameter for complex arithmetic as the smallest positive model number $z$ such that for any $x$ in the range $[z, 1 / z]$ the following can be computed without any undue loss of accuracy, overflow, underflow or other error:

$$
\begin{aligned}
& -w \\
& 1 / w \\
& -1 / w \\
& \operatorname{SQRT}(w) \\
& \operatorname{LOG}(w) \\
& \operatorname{EXP}(\operatorname{LOG}(w)) \\
& y * *(\operatorname{LOG}(w) / \operatorname{LOG}(y)) \text { for any complex } y \\
& \operatorname{ABS}(w)
\end{aligned}
$$

where $w$ is any of $x, i x, x+i x, 1 / x, i / x, 1 / x+i / x$, and $i$ is the square root of -1 .
This parameter was introduced to take account of the quality of complex arithmetic on the machine. On machines with well implemented complex arithmetic, its value will differ from that of the real safe range parameter by a small multiplying factor less than 10 . For poorly implemented complex arithmetic this factor may be larger by many orders of magnitude.

### 2.2 Other Aspects of the Computing Environment

No attempt has been made to characterise comprehensively any other aspects of the computing environment. The other functions in this chapter provide specific information that is occasionally required by routines in the Library.

### 2.3 References

[1] Brown W S (1981) A simple but realistic model of floating-point computation ACM Trans. Math. Software 7 445-480
[2] IEEE (1985) Standard for Binary Floating Point Arithmetic volume Standard 754-1985 ANSE/IEEE, New York.

## 3 Recommendations on Choice and Use of Available Routines

Note. Refer to the Users' Note for your implementation to check that a routine is available.
Routines in this Chapter need not be preceded by a call to Z01AAFP.

### 3.1 Parameters of Floating-point Arithmetic

DOUBLE PRECISION FUNCTION XO2AJF()

DOUBLE PRECISION FUNCTION XO2AKF()
DOUBLE PRECISION FUNCTION XO2ALF()

DOUBLE PRECISION FUNCTION XO2AMF()

DOUBLE PRECISION FUNCTION XO2ANF()

INTEGER FUNCTION X02BHF()
INTEGER FUNCTION X02BJF()
INTEGER FUNCTION XO2BKF()
INTEGER FUNCTION XO2BLF()
LOGICAL FUNCTION XO2DJF()
returns the machine precision, i.e., $\frac{1}{2} b^{1-p}$ if ROUNDS is true or $b^{1-p}$ otherwise (or a value very slightly larger than this, see Section 2.1.2)
returns the smallest positive model number, i.e., $b^{e_{\min }-1}$ returns the largest positive model number, i.e., ( $1-$ $\left.b^{-p}\right) b^{e_{\text {max }}}$
returns the safe range parameter as defined in Section 2.1.2
returns the safe range parameter for complex arithmetic as defined in Section 2.1.2
returns the model parameter $b$ returns the model parameter $p$
returns the model parameter $e_{\text {min }}$
returns the model parameter $e_{\max }$ returns the model parameter ROUNDS

### 3.2 Parameters of Other Aspects of the Computing Environment

DOUBLE PRECISION FUNCTION XO2AHF (X) DOUBLE PRECISION X

INTEGER FUNCTION X02BBF()
INTEGER FUNCTION XO2BEF()

INTEGER FUNCTION XO2DAF()
DOUBLE PRECISION X
returns the largest positive DOUBLE PRECISION argument for which the sin and cos routines return a result with some meaningful accuracy. X is a dummy argument
returns the largest positive integer value returns the maximum number of decimal digits which can be accurately represented over the whole range of floating-point numbers returns FALSE if the system sets underflowing quantities to zero, without any error indication or undesirable warning or system overhead. X is a dummy argument.

## 4 Example

The Example Program listed below simply prints the values of all the functions in Chapter X02. Obviously the results will vary from one implementation of the Library to another. The results listed in Section 4.3 are those from the Silicon Graphics (IRIX 4) implementation.

### 4.1 Example Text

```
* X02AJF Example Program Text
* NAG Parallel Library Release 2. NAG Copyright }199
* .. Parameters ..
    INTEGER NOUT
    PARAMETER (NOUT=6)
* .. External Functions ..
    DOUBLE PRECISION X02AHF, X02AJF, X02AKF, X02ALF, X02AMF, X02ANF
    INTEGER X02BBF, X02BEF, X02BHF, X02BJF, X02BKF, X02BLF
    LOGICAL X02DAF, X02DJF
    EXTERNAL X02AHF, X02AJF, X02AKF, X02ALF, X02AMF, X02ANF,
    + X02BBF, X02BEF, X02BHF, X02BJF, X02BKF, X02BLF,
    + X02DAF, X02DJF
* .. Executable Statements ..
    WRITE (NOUT,*) 'X02AJF Example Program Results'
    WRITE (NOUT,*)
    WRITE (NOUT,*) '(results are machine-dependent)'
    WRITE (NOUT,*)
    WRITE (NOUT,*) 'The basic parameters of the model'
    WRITE (NOUT,*)
    WRITE (NOUT,99999) ' X02BHF = ', X02BHF(),
    + ' (the model parameter B)'
    WRITE (NOUT,99999) ' X02BJF = ', X02BJF(),
    + , (the model parameter P)'
        WRITE (NOUT,99999) ' X02BKF = ', X02BKF(),
        + ' (the model parameter EMIN)'
        WRITE (NOUT,99999) ' XO2BLF = ', XO2BLF(),
    + , (the model parameter EMAX)'
        WRITE (NOUT,99998) ' X02DJF = ', X02DJF(),
    + , (the model parameter ROUNDS)'
        WRITE (NOUT,*)
        WRITE (NOUT,*)
        + 'Derived parameters of floating-point arithmetic'
        WRITE (NOUT,*)
        WRITE (NOUT,*) ' X02AJF = ', X02AJF(),
    + ' (the machine precision)'
```

```
        WRITE (NOUT,*) ' X02AKF = ', X02AKF(),
        + , (the smallest positive model number)'
        WRITE (NOUT,*) ' X02ALF = ', X02ALF(),
    + ' (the largest positive model number)'
        WRITE (NOUT,*) ' X02AMF = ', X02AMF(),
        + ' (the real safe range parameter)'
        WRITE (NOUT,*) ' X02ANF = ', X02ANF(),
        + , (the complex safe range parameter)'
        WRITE (NOUT,*)
        WRITE (NOUT,*)
    + 'Parameters of other aspects of the computing environment'
        WRITE (NOUT,*)
        WRITE (NOUT,*), X02AHF = ', XO2AHF(0.0DO),
    + ' (largest argument for SIN and COS)'
        WRITE (NOUT,99997)' X02BBF = ', X02BBF(0.0D0),
    + ' (largest positive integer)'
        WRITE (NOUT,99997) ' X02BEF = ', X02BEF(0.0DO),
    + , (precision in decimal digits)'
        WRITE (NOUT,99996) ' X02DAF = ', XO2DAF(0.0DO),
    + , (indicates how underflow is handled)'
        STOP
*
99999 FORMAT (1X,A,I7,A)
99998 FORMAT (1X,A,L7,A)
99997 FORMAT (1X,A,I20,A)
99996 FORMAT (1X,A,L20,A)
        END
```


### 4.2 Example Data

None.

### 4.3 Example Results

X02AJF Example Program Results
(results are machine-dependent)

The basic parameters of the model

```
X02BHF = 2 (the model parameter B)
X02BJF = 53 (the model parameter P)
X02BKF = -1021 (the model parameter EMIN)
X02BLF = 1024 (the model parameter EMAX)
X02DJF = T (the model parameter ROUNDS)
```

Derived parameters of floating-point arithmetic

```
X02AJF = 1.1102230246251600E-16 (the machine precision)
XO2AKF = 2.2250738585072107-308 (the smallest positive model number)
X02ALF = 1.7976931348623093+308 (the largest positive model number)
X02AMF = 2.2250738585072107-308 (the real safe range parameter)
X02ANF = 2.2250738585072107-308 (the complex safe range parameter)
```

Parameters of other aspects of the computing environment
$\mathrm{X} 02 \mathrm{AHF}=1.8014398509481900 \mathrm{E}+16$ (largest argument for SIN and COS)

```
X02BBF =
    2147483647 (largest positive integer)
X02BEF =
X02DAF = F (indicates how underflow is handled)
    15 (precision in decimal digits)
```

